



SF-8329

B. E. - III (Sem. - VI) Examination

May/June - 2011

Theory of Computation

Time : 3 Hours]

[Total Marks : 100

Instruction :

नीचे दशांश देव निशानीवाणी विगतो उत्तरवडी पर अवश्य लपवी.  
 Fillup strictly the details of signs on your answer book.

Name of the Examination :  
 B. E. - 3 (SEM. - 6)

Name of the Subject :  
 THEORY OF COMPUTATION

Subject Code No. : 8 3 2 9 Section No. (1, 2,.....): Nil

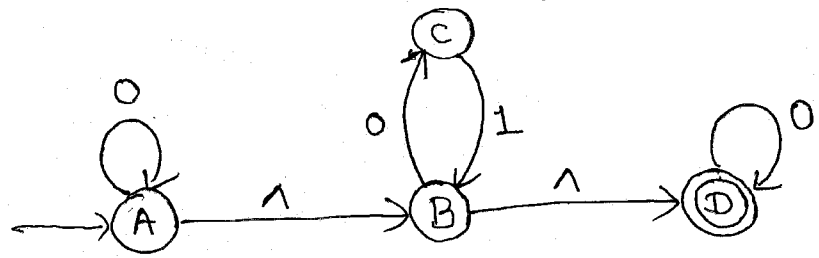
Seat No. :

Student's Signature

- 1 [A] Do as directed. [10]
- (1) Define one-to-one and on-to function. [02]
  - (2) Give recursive definition for set of all strings in  $\{0,1\}^*$  containing substring 00. [04]
  - (3) Explain the logical quantifiers and quantified statement. [02]
  - (4) Define FA. [02]

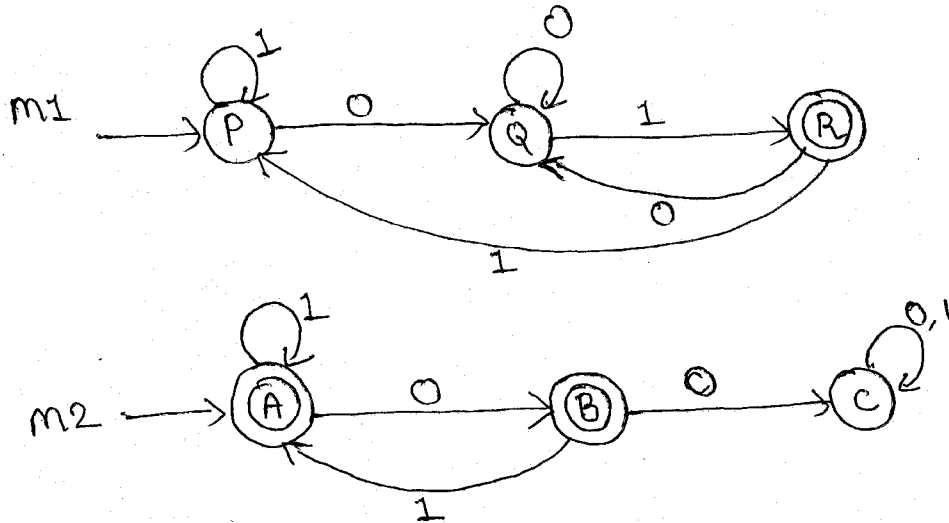
- 1[B] Give the regular expression for the following set of languages over set  $\{0,1\}$ . [04]
- (i) The language of all strings containing substring 110.
  - (ii) The language of all strings containing exactly two 0's.

- 2 Do as directed.
- (1) Draw an FA to recognize the following languages defined over  $\{0,1\}^*$ . (Any two)- [08]
    - (i)  $0 + 10^* + 01^*0$
    - (ii)  $(111 + 100)^* 0$
    - (iii)  $\{w \mid w \text{ has two or three } 1\text{'s}\}$
  - (2) Find equivalent NFA and FA for NFA -  $\Lambda$  given below. [08]



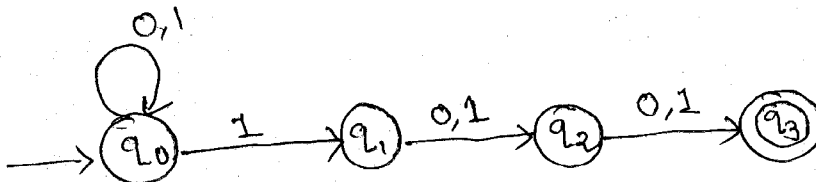
OR

- (2) Let  $M_1$  and  $M_2$  be the FA shown in figures, recognizing language  $L_1$  and  $L_2$  respectively. Draw an FA recognizing following languages. [08]
- $L_1 \cup L_2$
  - $L_1 \cap L_2$
  - $L_1 - L_2$



3 Do as directed. (Any four) [20]

- Define NFA -  $\Lambda$  and extended transition function  $\delta^*$  for NFA -  $\Lambda$ .
- Using pumping lemma prove that the language  $L = \{0^i 1^j 0^k \mid k > i + j\}$  is not regular.
- Draw an NFA -  $\Lambda$  recognizing the language  $((ab)^* b + ab^* )^*$ .
- Calculate  $\delta^*(q_0, 01)$ ,  $\delta^*(q_0, 011)$  for NFA given below.



- (5) Describe P, NP and NP complete problem..

4 [A] Do as Directed [10]

- Define Semi word
- Define Context Free Grammar (CFG)
- The language EVENPALINDROME is a regular language (TRUE/FALSE)
- PDA maintains a tape for storage. (TRUE/FALSE)
- The intersection of two context free language is always context free (TRUE/FALSE)
- The class of automata that can be associated with context free languages is \_\_\_\_\_.
- A CFG  $G$  is said to be \_\_\_\_\_ if there exists a string  $w$  in  $L(G)$ , for which more than one parse trees can be generated.

- (8) Recursively Enumerable languages are accepted by \_\_\_\_\_.
- (9) Turing machine uses \_\_\_\_\_ as storage device.
- (10) The set of  $\{a^n b^{2n} \mid n = 1, 2, \dots\}$  can be generated by the CFG \_\_\_\_\_.

[B] Define **Chomsky Normal Form (CNF)**. [10]

Convert the following CFG to CNF

- (i)  $E \rightarrow E+E \mid E^*E \mid (E) \mid \epsilon$
- (ii)  $S \rightarrow SaS \mid SaSbS \mid SbSaS \mid \epsilon$

5 [A] Find **CFG** for the following languages [8]

(i)  $\{ a^i b^j c^k \mid i, j, k = 1, 2, 3, \dots \text{ And } i \neq j \text{ or } j \neq k \}$

(ii)  $\{ a^i b^j c^k \mid i, j, k = 1, 2, 3, \dots \text{ And } j = 2i + 2k \}$

[B] Build **Turing Machine** to compute 1's complement. [8]

**OR**

[A] Give the **CFG** for the language  $a^n b^n$ . Convert the CFG to CNF. Make PDA that accepts the language. [8]

[B] Build **Turing Machine** for subprogram **DELETE**. [8]

6 [A] Attempt **Any one**. [6]

(1) Design a Turing Machine to recognize the string of the type  $(01)^n$ .

(2) For the following CFG's describe the language it accepts:

(a)  $S \rightarrow SS \mid XaXaX \mid \epsilon$

$X \rightarrow bX \mid \epsilon$

(b)  $S \rightarrow aS \mid bS \mid a \mid b \mid \epsilon$

(c)  $S \rightarrow aM \mid bS$

$M \rightarrow aF \mid bS$

$F \rightarrow aF \mid bF \mid \epsilon$

[B] Attempt **Any Two**. [8]

(1) Prove that "The Context Free Languages are closed under union."

(2) Killing the Unit Productions

$S \rightarrow A \mid bb$

$A \rightarrow B \mid b$

$B \rightarrow S \mid a$

(3) Design the CFG for following

Regular Expression =  $01(1100)^*110(10)^*$ ,

Set of NonTerminals =  $\{S, A, B\}$ ,  $\Sigma = \{0, 1\}$